

Faster Point Doubling on Twisted Hessian Curves

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Short Weierstrass Curve

$$y^2 = x^3 + ax + b$$

where $4a^3 + 27b^2 \neq 0$

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$$y^2 = x^3 - 0.4x + 0.7$$



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Edwards Curve

$$x^2 + y^2 = c^2(1 + dx^2y^2)$$

where $c \neq 0$ and $d \notin \{0, 1\}$

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$$x^2 + y^2 = 1 - 300x^2y^2$$



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Hessian Curve

$$x^3 + y^3 + 1 = 3dxy$$

where $d \neq 1$

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$$x^3 - y^3 + 1 = 0.3xy$$



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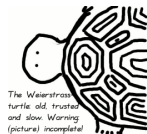


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Curve Shapes

- Short Weierstrass curves

$$y^2 = x^3 - 0.4x + 7$$



- Edwards curves

$$x^2 + y^2 = 1 - 300x^2y^2$$



- Hessian curves

$$x^3 - y^3 + 1 = 0.3xy$$

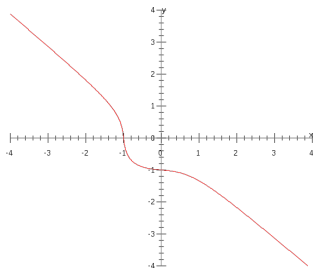


<https://hyperelliptic.org/EFD/>

Twisted Hessian Curves

- Hessian curves:

$$x^3 + y^3 + 1 = 3dxy$$

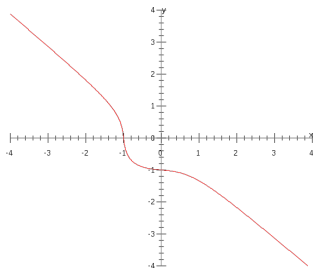


$$x^3 + y^3 + 1 = 0.3xy$$

pictures credit http://en.wikipedia.org/wiki/Hessian_form_of_elliptic_curve

Twisted Hessian Curves

- Hessian curves:
 $x^3 + y^3 + 1 = 3dxy$
- Twisted Hessian curves:
 $ax^3 + y^3 + 1 = 3dxy$



$$x^3 + y^3 + 1 = 0.3xy$$

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Scalar Multiplication

- Given scalar n and point P
- Compute $nP = \underbrace{P + P + \dots + P}_n$

Scalar Multiplication

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- Algorithm: Double-and-add

Double-and-add Algorithm

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- E.g. $n = 23$

Double-and-add Algorithm

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- E.g. $n = 23 = 10111_2$
 $\quad \quad \quad P \quad \quad \quad : P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2P$

: $2P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2P)$

: $4P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2P) + P$: $5P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2(2P) + P) \quad : 10P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2(2P) + P) + P$: $11P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2(2(2P) + P) + P) \quad : 22P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$

- E.g. $n = 23 = 10111_2$
 $2(2(2(2P) + P) + P) + P \quad : 23P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$
 $23P = 2(2(2(2P) + P) + P) + P$

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$
 $23P = 2(2(2(2P) + P) + P) + P$
- Number of operations:
 - Double = # of bit - 1
 - Addition = # of bit set - 1

Double-and-add Algorithm

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$
 $23P = 2(2(2(2P) + P) + P) + P$
- Number of operations:
 - Double = # of bit - 1
 - Addition = # of bit - 1 (add always)

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$

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- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = \mathbf{1}0111_2$, $w = 2$
 $\quad \quad \quad P \quad \quad \quad : P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $2^2(P)$: $4P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $2^2(P) + P : 5P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $2^2(2^2(P) + P) : 20P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $2^2(2^2(P) + P) + 3P$: $23P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $23P = 2^2(2^2(P) + P) + 3P$

- Compute $nP = \underbrace{P + P + \dots + P}_n$
- E.g. $n = 23 = 10111_2$, $w = 2$
 $23P = 2^2(2^2(P) + P) + 3P$
- Number of operations:
 - Double = # of bit - 1
 - Addition = $\lceil (\# \text{ of bit} - 1) / w \rceil$

Previously best doubling formulas
(dbl-2009-bkl-3):

$$i_Z = i \cdot Z_1$$

$$A = (Y_1 - i_Z) \cdot (Y_1 + i_Z)$$

$$B = Y_1 \cdot Z_1$$

$$C = (A - B) \cdot (Y_1 + Z_1)$$

$$D = (A + B) \cdot (Z_1 - Y_1)$$

$$E = 3 \cdot C - 2d \cdot X_1 \cdot B$$

$$X_3 = (-2) \cdot X_1 \cdot D$$

$$Y_3 = (D - E) \cdot Z_1$$

$$Z_3 = (D + E) \cdot Y_1$$

Faster formulas

Previously best doubling formulas
(dbl-2009-bkl-3):

$$i_Z = i \cdot Z_1$$

$$A = (Y_1 - i_Z) \cdot (Y_1 + i_Z)$$

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$$E = 3 \cdot C - 2d \cdot X_1 \cdot B$$

$$X_3 = (-2) \cdot X_1 \cdot D$$

$$Y_3 = (D - E) \cdot Z_1$$

$$Z_3 = (D + E) \cdot Y_1$$

New formulas:

$$P = Y_1 \cdot Z_1$$

$$2P = 2 \cdot P$$

$$S = Y_1 + Z_1$$

$$A = S^2 - P$$

$$C = (A - 2P) \cdot S$$

$$D = A \cdot (Z_1 - Y_1)$$

$$E = 3 \cdot C - d \cdot X_1 \cdot 2P$$

$$X_3 = (-2) \cdot X_1 \cdot D$$

$$Y_3 = (D - E) \cdot Z_1$$

$$Z_3 = (D + E) \cdot Y_1$$

dbl-2009-bkl-3 fomulas:

- Assumption:
 - $i^2 = -1$
 - $2d = 2 \cdot d$

- Cost

8M

+ 1*(-2)

+ 1*2d + 9*add

+ 1*3

+ 1*i

New formulas:

- Cost

7M + 1S

+ 1*(-2)

+ 1*d + 7*add

+ 1*3

+ 1*2